Preliminary Study on Optimization of Data Distribution in Resource Sharing Systems

Grzegorz Chmaj, Krzysztof Walkowiak
Chair of Systems and Computer Networks, Wroclaw University of Technology, Poland
Grzegorz@Chmaj.net, Krzysztof.Walkowiak@pwr.wroc.pl

Abstract

Grid structures are increasingly considered as very convergent with peer-to-peer networks. This paper presents a model of network acting both as grid and peer-to-peer network, used for data computation and distribution. Presented PPLC algorithm is a complete solution for both grid and peer-to-peer points of view. Problem formulation is presented, as well as solution heuristic algorithm and research results.

1. Introduction

Peer-to-peer (p2p) networks are structures containing many nodes having the same privileges and performing the same tasks [2], [4]. The one paradigm of p2p network is that each node has the same role in such network (server and client). However, some peer-to-peer systems (e.g. BitTorrent) use many roles – some nodes may perform additional management tasks. Grid networks are built to achieve processing power which is spread across computation centers. Each grid participant offers computation resources – the computation problem is divided into fragments and each fragment is independently processed. Then results can be combined back into one product [3], [6]. By source data we mean the input data which requires large computation power to be processed, so it cannot be done at one computing unit. Most known grid projects are Seti@home (searching for extra terrestrial intelligence), Folding@home (protein folding and other molecular dynamics). Many components were created to support grid development: APIs (OGSA, OGSI, SAGA, CORBA, DRMAA, GSI) and software implementations (BOINC, Globus Toolkit).

Grids and p2p are converged in many aspects [1], one common architecture could be used (or interpreted) as grid or p2p network [5]. This paper describes joint of grid and peer-to-peer network, used to compute data and distribute results to all participants. Initial (“source”) data is divided into blocks and distributed among nodes of network (blocks represent network data units). Network – acting as grid structure – performs computation of blocks. Each node process blocks it has assigned to itself, computation cost may differ among nodes. To become project's participant, such node has to process at least one source data block. Computation process transforms blocks into new form, considered as “result” data. Each node computes blocks autonomously, without interaction with other nodes. Next, result data is distributed to all participating nodes (as all blocks have to be distributed to all nodes) using peer-to-peer mechanisms. We model peer-to-peer transfer as set of iterations which may be considered as time slots. Each of iteration allows performing some actions, with results visible in the next time slot. Both computation of source block and data transfer between two nodes of network introduces a cost which is the objective to be minimized. Whole process has also time limit, determined by iterations.

We present a network model covering all above aspects, suitable for systems performing data processing and result sharing. Both optimal and heuristic solutions with results are presented. This kind of system can be applied for example to geographic maps processing performed by set of collaborators, which need to have all results visible locally. In that case we would have the map structure stored in each processing cluster. Map would be divided into fragments independently processed by various units.

This paper is organized as follows: section 2 introduces nomenclature used in further considerations. Section 3 contains problem formulation. PPLC algorithm is described in section 4, results of experiments are showed in section 5. Conclusions are formulated in section 6.
2. Nomenclature

Following terms will be used to comprehensively describe the problem and its solution: block – represents data fragment denoted as b, that can be processed on network node and sent between network nodes; vertex (network node) – denoted as v, w, s - element of the network, that is able to process data blocks, send them and fetch to/from other nodes; network – set of V nodes connected with each other. Vertices that belong to one network may share information between each other. iteration (time slot) – denoted as i. We consider network processing as the set of iterations. In each of iteration, nodes may transfer blocks between them, but information about assignment of blocks to nodes is updated at the beginning of next iteration. More ways of network modeling you can find in [7]; processing – block b can be computed (processed) on network node. Resulting data replaces original (source) data. Block b can be processed only by vertex which has block b assigned to; block to node assignment – block b is assigned to node, when it is stored physically on given node v – then, all other nodes may fetch block b from node v; transfer – the process of sending data (blocks) from source node to destination node (then block transferred to node v becomes assigned to node v): peer-to-peer network (abbreviated to p2p network) – network in which all nodes act both as clients and servers [2]; grid – set of nodes, considered as one body that is able to perform common tasks; saturated network – network is saturated, when every node in the network has all blocks assigned to itself (4).

3. Problem formulation

To mathematically represent the problem we introduce the following notations:

indices
\[ b = 1, \ldots, B \quad \text{blocks to be transferred} \]
\[ t, i = 1, \ldots, T \quad \text{time slots (iterations)} \]
\[ v, w, s = 1, \ldots, V \quad \text{vertices (network nodes)} \]

constants
\[ c_v \quad \text{cost of processing block in node } v \]
\[ k_{bw} \quad \text{cost of block transfer from node } w \text{ to node } v \]
\[ M \quad \text{large number} \]

variables
\[ x_{bv} = 1 \text{ if block with index } b \text{ is processed (calculated) in node } v; 0 \text{ otherwise (binary)} \]
\[ y_{bwv} = 1 \text{ if block } b \text{ is transferred to node } v \text{ from node } w \text{ in iteration } t; 0 \text{ otherwise (binary)} \]

objective
\[ \min: F = \sum \sum x_{bv} c_v + \sum \sum \sum y_{bwv} k_{bw} \] (1)

constraints
\[ \sum_b x_{bv} \geq 1 \quad v = 1, \ldots, V \] (2)
\[ \sum_v x_{bv} = 1 \quad b = 1, \ldots, B \] (3)
\[ x_{bv} + \sum_v y_{bwv} = 1 \quad b = 1, \ldots, B \quad v = 1, \ldots, V \] (4)
\[ \sum_v y_{bwv} \leq M (x_{bw} + \sum_i \sum_v y_{bwv}) \quad b = 1, \ldots, B \] (5)

The objective function (1) is the cost of processing of all blocks (the first term) and the cost of blocks’ transfer (the second term). Condition (2) states that each vertex must process at least one block. Constraint (3) assures that each block is assigned to only one network vertex. To meet the requirement that each processed block must be transported to each network node we introduce the condition (4), (5) defines the source node from which block b may be transferred. M variable should be larger than VB to guarantee that block b resides on node w.

4. PPLC Algorithm

In this section we describe a heuristic algorithm PPLC (Pre-Post Limited Cost) that solves the problem (1-5). The PPLC algorithm (see Fig. 1) consists of two phases. The first phase is aimed to assign blocks to vertices for processing, i.e. values of variables \( x_{bv} \) are selected. The second part of PPLC is responsible for the P2P-based block transfer, i.e. values of variables \( y_{bwv} \) are found. The overall objective of PPLC is to find a solution that satisfies constraints (2-5) and minimizes the cost function given by (1). Whole processing has to be completed in T iterations.

4.1. Assignment of blocks

In this subsection we will present how the PPLC algorithm assigns blocks to nodes for computation. General allocation schema is presented on Fig. 2.

Notice that the block b can be assigned to node v for processing (variable \( x_{bv}=1 \)) or block b is transferred to node v in one of iterations (variable \( y_{bwv}=1 \)). At the beginning one block is allocated to every node – to satisfy formula (2) (V blocks are allocated). Then, if \( V < B \), PPLC performs II level allocation. For every node \( v=1, \ldots, V \) score is calculated using formula (6), and maximum value among score values is chosen (7). For every node \( v=1, \ldots, V \) score gap \( g_v \) is computed (8), and all \( g_v \) values are put into a G array which is then
allocate blocks to vertices, according to (2) and (3)
2 process blocks’ computation
3 compute m value.
4 find sets of nodes having distance ≥ m (CVGs). If found more than 1 CVGs, then reserve first iteration for TR0 and transport blocks between all CVGs (TR0)
8 if f = w and f = v and exist at least one node v which transport costs to all other nodes are greater or equal m, then reserve last iteration for TR3
9 perform first main (not reserved) iteration: compute f, perform all possible transfers having cost ≤ f,
10 perform TR1 phase (see Fig. 3)
11 for all not-reserved iterations i: compute f_i and perform all possible transfers having cost less or equal f_i
12 after all transfers in last main iteration perform TR2 (see Fig. 4)
13 if last iteration is not reserved, perform TR3 phase, otherwise go to last iteration and then perform TR3 phase

Fig. 1 Pseudocode of PPLC algorithm
1 allocate one block to each node
2 if all blocks are allocated – end of allocation. Otherwise go to step 3.
3 compute score w for each node v
4 compute score gap g, for each node v
5 allocate blocks to nodes having highest g, unless all blocks are allocated

Fig. 2 Pseudocode of blocks’ allocation

sorted by value descending. Let us notice, that element g of array G is the value of gap (g_v) and we also have the information to which node this particular gap value is assigned. Let first(G) return the first element of G.
G=G−{g} denotes deletion of element g from array G, so then next element becomes the first one.

\[ e_v = c_v + \sum k_{vw} \]  
\[ e_{max} = \max (e_v) \]  
\[ g_v = \frac{e_v - e_{max}}{e_v} \]  

Having array of g values, PPLC allocates blocks to nodes using following schema: Let \( a_{sum} \) be the number of blocks allocated in II level of allocation, and \( a_{sum} = 0 \). Allocate \( a_v \), (9) blocks to node associated with g.
\[ a_v = \begin{cases} g \cdot (B - V) & \text{if } (B - V)(1 - g) - a_{sum} \geq 0 \\ B - a_{sum} & \text{otherwise} \end{cases} \]  

While there are still unallocated blocks, set \( G = G - \{g\} \), \( v = \text{first}(G) \) and \( a_{sum} \) is increased by \( a_v \) and allocation is performed again based on (9). This process is repeated unless all blocks are allocated. When all blocks are assigned to nodes, computation process is performed.

4.2. Transfer of blocks

In this subsection we will explain how the PPLC algorithm distributes the blocks to all vertices. Recall that according to our model (1-5) the P2P approach is applied for distribution of blocks. First, we introduce parameter \( m \) defined as follows
\[ m = \min (k_{vw}) \ v, w = 1, ..., V \ v \neq w \]  
\[ m = \max (m_v) \ v = 1, ..., V \]  

Value \( m \) (11) is the coefficient used to scale cost limit. It has constant value during processing and is established at the beginning of algorithm, acting as general constant parameter.

PPLC limits the maximal cost in each operation \( i \) to settled value \( f_i \) (12), which is computed for each of iteration separately.
\[ f_i = \left[ \frac{m_i}{T} \right] \]  

High values of \( m \) would let blocks to be transferred at higher costs, what is not favorable according to objective (1). Too low value of \( m \) could cause such situation, that many blocks are transferred in TR3 phase (what implies high costs). Value of \( m \) is the maximum cost of acceptable transfer during iterations, except TR0 - and TR3.

Now we define CVG (closed vertices group). A CVG contains all vertices that are within distance limited by m. More precisely, any two nodes \( w \) and \( v \) belong to the same CVG denoted as \( Q \) if \( k_{vw} \leq m \). If for any two nodes \( w \) and \( v \) the following condition holds \( k_{vw} > m \), it means that \( w \) and \( v \) do not belong to the same CVG. If (13) and (14) are satisfied – then \( Q_i \) constitutes CVG.
\[ k_{vw} \leq m \ v, w \in Q_i \]  
\[ k_{vw} > m \ v, w \notin Q_i \]  

The PPLC algorithm finds all CVGs. When more than one such group is found then the first iteration is performed only for blocks’ transport between CVGs, which is performed in the following way. Firstly, PPLC creates a set \( K \) of containing pairs of indices identifying CVGs. All possible combinations of CVG pairs are put into \( K \). Value of \( f_i \) is computed using the following formula (15)
\[ f_i = \max_k k_{vw} \ v, w = 1, ..., V \]  

We assume that the operator first(K) returns the first pair of CVGs from \( K \). Let \( (Q_1, Q_2) = \text{first}(K) \).

Let \( D_i \) denote the set of blocks which are allocated to node \( v \) (\( D_n \) and \( D_t \) – by analogy). PPLC performs transfer from \( Q_1 \) to \( Q_2 \) (16) and from \( Q_2 \) to \( Q_1 \) (17).
\[ y_{bwv} = 1 \ v \in Q_1 \ w \in Q_2 \ b \in D_i \ t = 1 \]  
\[ y_{bwv} = 1 \ w \in Q_2 \ v \in Q_1 \ b \in D_o \ t = 1 \]  

Destination nodes: \( w \) for transfer \( Q_1 \) to \( Q_2 \), and \( v \) for transfer \( Q_2 \) to \( Q_1 \) should satisfy conditions (18) (for node \( v \)) and (19) (for node \( w \)).
\[ k_{vw} \leq m \ v \in Q_1 \ s \in \overline{Q_1} \]  
\[ k_{vw} \leq m \ w \in Q_2 \ s \in \overline{Q_2} \]
1 compute \( f_{rm} \);
2 for every node \( v=1,...,V \) do:
3 if transport cost from \( s \) is less or equal \( m \) to less than 3 nodes
   or: transport cost from \( s \) to another node \( w \) is less or equal \( f_{rw} \) and transport cost from \( s \) to all nodes other
   than \( w \) is greater than \( f_{rw} \), then: transfer all block available on \( s \) to node \( v \)

**Fig. 3 Pseudocode of TR1**

1 for every node \( v=1,...,V \) do:
2 if transport cost from \( v \) is less or equal \( m \) to less than 3 nodes then find node \( s \) which has smallest transfer cost to node \( v \)
3 for every block \( b \) absent on \( v \) and absent on \( s \):
4 \( tcost_i=\) cost of sending \( b \) to \( v \) in current iteration
5 \( tcost_i=\) cost of sending \( b \) to \( s \) in current iteration
6 \( tcost_i=\) cost of sending \( b \) to \( v \) in last iteration
7 \( tcost_i=\) cost of sending \( b \) to \( s \) in last iteration
8 if \( tcost_{1}+tcost_{2}+tcost_{3} \) then transfer \( b \) to \( s \)

**Fig. 4 Pseudocode of TR2**

If these conditions cannot be satisfied due to network topology, nodes \( v \) and \( w \) are chosen in a way: that for transport \( Q_1 \) to \( Q_2 \), node \( v \) is the one that has the biggest quantity of transfer costs to other nodes belonging to \( Q_2 \); \( Q_1 \) to \( Q_3 \), node \( w \) is the one that has most quantity of transfer costs to other nodes belonging to \( Q_1 \).

Transfers between CVGs are performed for all elements from \( K \) array, so transfers between all possible pairs of CVGs are done. If there were more than one CVG groups found, then first operation is purposed only for TR0. At this point, PPLC decides if last iteration should be purposed for transfers that cost more than \( m \) (iteration \( =T \) is purposed for TR3 phase only). If \( f_{rm}=m \) and \( f_{rw}=m \) (12) and there exists at least one node \( v \) having all transfers costs to other nodes greater or equal \( m \) (20), then last iteration is reserved for over \( m \) cost transfers. Let us name not-reserved iterations as main iterations. The next step is to transfer blocks using all iterations except reserved earlier.

\[
k_{vw} \geq m \quad w=1,...,V \quad w \neq v
\]  

(20)

For each iteration \( i \) (except iterations reserved for TR0 and TR3 if such reservation was made), value of \( f_i \) is computed using (12). Then, all possible transfers having cost less or equal \( m \) are performed (21) – nodes fetch only blocks, which they do not possess yet.

\[
y_{bwv}=1 \text{ for } b=1,...,B \quad w=1,...,V, \quad w \neq v, \quad k_{vw} \leq f_i \text{ and } x_{bw} + \sum_y y_{bwv} = 0
\]  

(21)

If this is first main iteration, then TR1 transfer is performed, according to schema presented on Fig. 3.

First step of TR1 is to compute the cost limit for iteration \( T-1 \) (12). PPLC finds the most attractive node \( v \) – node having the biggest quantity of transfer costs less or equal \( m \) (24).

\[
h_{vs} = \begin{cases} 
1 & \text{if } k_{vs} \leq m \\
0 & \text{otherwise}
\end{cases}
\]  

(22)

\[
H_{vs} = \sum y_{swv}
\]  

\[
v = \max_{s} H_{vs} \quad s = 1,...,V
\]  

(23)

(24)

Then, for each node \( s \) following checks are made: a) if number of network connections having transfer costs less or equal \( m \) from node \( s \) is less than 3; b) node \( s \) has the following transport costs’ structure: there is one network node, to which \( s \) has transport cost less or equal \( m \), and transport cost to all other nodes is greater than \( m \). If a) or b) is true, then all blocks existing on \( s \) and not existing on node \( v \) are transferred from \( s \) to \( v \). After TR1 phase, PPLC algorithm continues with main iterations. At the end of last main iteration, phase TR2 is performed, according to Fig. 4: for each node \( v \), following steps are performed: if number of network connections having transfer costs less or equal \( m \) from node \( v \) is less than 3, then node \( s \) is found. Node \( s \) is selected based on criterion: node \( s \) is the one that has minimal transfer cost to node \( v \). Then, for every block \( b \) that does not exist on \( v \) (25) and does not exist on node \( s \) (26): if \( t_{cost_1} + t_{cost_2} + t_{cost_3} + t_{cost_4} \) (please see (27), (28), (29), (30)) then PPLC transfers block \( b \) to node \( s \) from node having smallest transfer cost to \( s \).

\[
x_{bw} + \sum_y y_{bwv} = 0
\]  

(25)

\[
x_{bw} + \sum_y y_{bwv} = 0
\]  

(26)

\[
t_{cost_1} = k_{vw} \text{ where } w = \min_{w} k_{vw}
\]  

(27)

\[
\text{and } x_{bw} + \sum_y y_{bwv} = 0
\]  

\[
t_{cost_2} = k_{vw} \text{ where } w = \min_{w} k_{vw}
\]  

(28)

\[
\text{and } x_{bw} + \sum_y y_{bwv} = 0
\]  

\[
t_{cost_3} = k_{vw} \text{ where } w = \min_{w} k_{vw}
\]  

(29)

\[
\text{and } x_{bw} + \sum_y y_{bwv} = 0
\]  

\[
t_{cost_4} = k_{vw} \text{ where } w = \min_{w} k_{vw}
\]  

(30)

The last step of PPLC is to transfer all blocks which cause network to be not saturated (TR3). This process is performed at the end of last iteration. During this transfer \( f_i \) is set basing on (15) and transfers are performed according to (21). At this point, all blocks are possessed by all nodes (4).

5. Experimentation results

Experimentical results of PPLC algorithm showed, that it is able to provide solution to the problem. Experimentation system was created using C++ under g++ compiler, and then compiled under Microsoft Visual Studio 2003. The system implements PPLC algorithm and is able to do value overriding, so values of coefficients may be set manually (this was used to experiments showed on Fig. 3). Network data generator
was also created – it produces network with parameters enclosed in given ranges.

For researches about PPLC algorithm itself, there were 20 networks generated – using following parameters given to network data generator: quantity of nodes: 50-90, quantity of blocks: 1-10 per node, transfer costs: 1-50. TR0 and TR3 affect on total transfer costs in networks having CVGs lowering total transport cost. For experiment where TR0 is not performed, not all blocks are transferred (transfers occurs only inside CVGs). Condition (4) is not satisfied then, so result where both TR0 and TR3 are not used cannot be accepted. For networks with CVGs when using only TR3, cost is highest – blocks between CVGs are transferred only during TR3 phase (what implies high costs). The relation between cost and m coefficient is presented on Fig. 5. To obtain this relation, experimental system was modified: value of m was not set using (11b), but it was set to fixed value. Experiments for given network were made for each value of m in range [1, 20]. Fig. 5 shows, that value of m should be low, to limit the cost in adequate way. Too low values cause cost to increase very quickly – more transfer is performed in TR3 on higher (> m) costs. Cost values have one minimum at \(m=m_0\), for values higher than \(m_0\), cost of transfer related to m is growing up. Lowering m values causes more blocks to be transferred in TR3, what is not efficient. Low values of m also result in presence of CVGs – and some transfer costs emerge in a phase of TR3. When values of m become higher, more blocks are transferred in main iterations – obeying \(f_i\) costs. Consequently, TR3 and TR0 perform less transfers cost. PPLC algorithm was also experimented in the scope of m designation. For 20 experimented networks, \(m=m_0\) was designated in 16 cases.

PPLC algorithm results were compared with optimal solution generated by CPLEX solver. Time limit for CPLEX computation process was set to 3600 seconds, in many cases CPLEX did not find optimal solution during this time, and returned feasible solution. On Fig. 6 we can observe the relation of solution values for both PPLC and CPLEX, and Fig. 7 shows parameters of networks used for experiments. All CPLEX results on this graph are optimal, and PPLC result values are 1%-20% close to optimal solution. The same experiments were made for larger networks, and results are presented on Fig. 8 and Fig. 9. In this case, CPLEX delivered feasible solution, as it was not able to produce optimal solution in 3600 seconds time. For smaller networks, CPLEX feasible solutions are still lower than PPLC, but for larger networks PPLC produces better solutions.

Fig. 10 shows results for one network (8 nodes, 5 iterations), where number of blocks was increased. We observe that for 160 blocks processed, PPLC and CPLEX solutions are still close, but for larger number of blocks PPLC algorithm results in much better solutions. Experiments showed that CPLEX is able to deliver optimal solutions in 3600 seconds only for
6. Conclusion

Presented approach of P2P-based data distribution in grid systems provides many processing possibilities, but requires efficient algorithm to manage transfers and blocks’ distribution. The joint of grid and peer-to-peer point of view has been proposed as one system, serving as grid or p2p at one time. Algorithm PPLC – proposed in this paper – achieves satisfying efficiency results of transfer costs and planning, suitable both for grid and p2p. Future work is to extend network model with real limitations, such as computation capacity and network transfers limitations.

7. References